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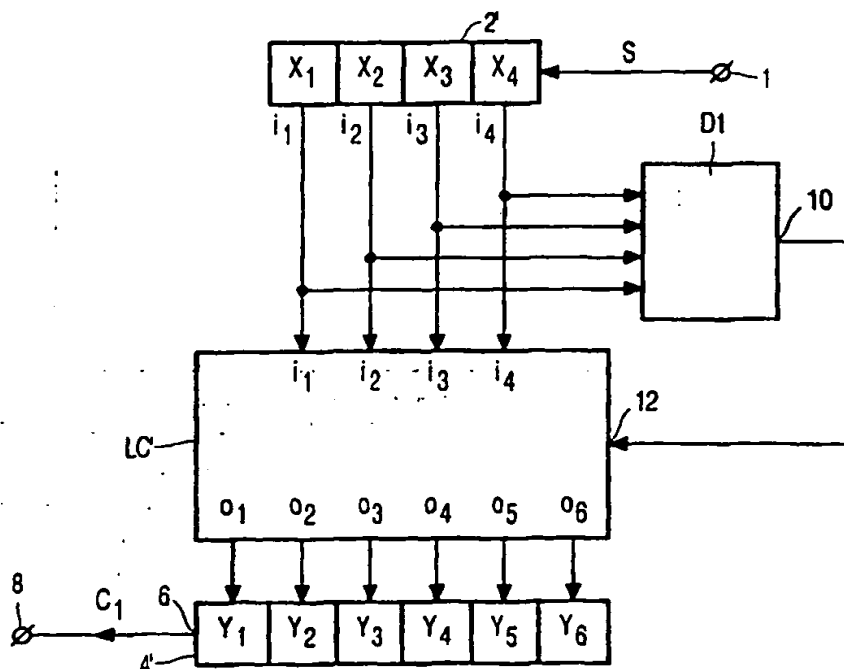
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(54) Title: ENCODING/DECODING n -BIT SOURCE WORDS INTO CORRESPONDING m -BIT CHANNEL WORDS, AND VICE
VERSA, SUCH THAT THE CONVERSION IS PARITY INVERTING

(57) Abstract

A device is disclosed for encoding a stream of databits of a binary source signal (S) into a stream of databits of a binary channel signal (C), wherein the bitstream of the source signal is divided into n -bit source words, which device comprises converting means adapted to convert said source words into corresponding m -bit channel words. The converting means are further adapted to convert a block of p consecutive n -bit source words into corresponding m -bit channel words, such that the conversion for each n -bit source word is parity inverting. The relations hold that $m > n \geq 1$, $p \geq 1$, and that p is an odd integer that can vary. Preferably, $m=n+1$. Further, a decoding device is disclosed for decoding the channel signal obtained by means of the encoding device.



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Encoding/decoding n-bit source words into corresponding m-bit channel words, and vice versa, such that the conversion is parity inverting.

The invention relates to a device for encoding a stream of databits of a binary source signal into a stream of databits of a binary channel signal, wherein the bitstream of the source signal is divided into n-bit source words, which device comprises converting means adapted to convert said source words into corresponding m-bit channel words and to a method.

5 The invention also relates to a device for decoding a stream of data bits of a binary channel signal obtained by means of the encoding device, so as to obtain a stream of databits of a binary source signal.

An encoding device mentioned in the foregoing is known from the book
10 'Coding techniques for digital recorders' by K.A. Schouhamer Immink, chapter 5.6.7, pp. 127 to 131, Prentice Hall (1991). The book discusses an encoder for generating a (d,k) sequence which satisfies the parameters: rate $2/3$, (1,7), which encoder is also proposed by Cohn et al in USP 4,337,458. The known encoding scheme suffers from the presence of a DC level which may become excessively large and therefore introduces distortion in communication systems
15 which can not handle a DC component, as well as distortion in any recording of data in magnetic media.

The invention has for its object to provide a device for encoding n-bit source words into corresponding m-bit channel words, such that it itself does not generate a DC
20 component in the channel signal, whereas further it provides the possibility, by means of additional measures to be taken, to realize a channel signal in the form of a (d,k) sequence.

The device in accordance with the invention is characterized in that the converting means are adapted to convert a block of p consecutive n-bit source words into a corresponding block of p consecutive m-bit channel words, such that the conversion for each
25 block of p consecutive n-bit source words is parity inverting, where n, m and p are integers, $m > n \geq 1$, $p \geq 1$, and where p is an odd integer and can vary.

'Parity inverting' means that the parity of the n-bit source words to be converted are the inverse of the parity (after modulo-2 addition) of the corresponding m-bit channel words in which they are converted. As a result, a unique relationship between the parity of the source words and

the parity of the channel words can be obtained, enabling an efficient DC control for the binary channel signal, after aT precoding.

The encoding device in accordance with the invention can be used in combination with a bit-adder unit in which one bit is added to codewords of a certain length.

5 The signal obtained can be applied to the encoding device of the present invention. The channel signal of the encoding device is applied to a 1T-precoder. The purpose of the bit-adder unit is to add a '0'- or a '1'-bit to blocks of data in the input signal of the converter, so as to obtain a precoder output signal which is DC free, or includes a tracking pilot signal having a certain frequency. The precoder output signal is recorded on a record carrier. The adding of a
10 '0'-bit in the input signal of the converter results in the polarity of the output signal of the 1T precoder remaining the same. The adding of a '1'-bit results in a polarity inversion in the output signal of the 1T precoder. The converter therefore influences the output signal of the 1T precoder such that the running digital sum value of the output signal of the 1T precoder can be controlled so as to have a desired pattern as a function of time.

15 Because of the fact that the encoding device in accordance with the invention realizes a parity-inverting encoding, it does not influence the polarity of the signal it encodes and can therefore be used in combination with the bit-adder unit without the need of any modification.

Preferably, m equals $n+1$, and n is equal to 2. For n being equal to 1 or 2, the
20 device in accordance with the invention can be used, with additional measures to be taken, as will become apparent later, for generating channel signals in the form of a (d,k) sequence, where $d=1$. Higher values for n do not allow for the generation of a $(1,k)$ sequence. Further, $n=1$, which means that 1-bit source words are converted into 2-bit channel words, results in an increase of 100 % in bits in the channel signal generated by the device. Contrary
25 to this a conversion of 2-bit source words into 3-bit channel words results in an increase of only 50 %, and is therefore more advantageous.

Various conversions of 2-bit source words into 3-bit channel words are possible, that have the parity inverting character. One such conversions is the subject of claim 4. It should however be noted that various permutations of the channel codes in the table are
30 possible, namely in total 4.

The device in accordance with the invention wherein the converting means are adapted to convert 2-bit source words into corresponding 3-bit channel words, so as to obtain a channel signal in the form of a (d,k) sequence, where $d=1$, the device further comprising means for detecting the position in the bitstream of the source signal where encoding of single

2-bit source words into corresponding single channel words would lead to a violation of the d-constraint at the channel word boundaries and for supplying a control signal in response to said detection, may be further characterized in that in the absence of the control signal, the converting means are adapted to convert single 2-bit source words into corresponding single 3-bit channel words, such that the conversion for each 2-bit source word is parity inverting.

More specifically, the device is characterized in that, in the presence of said control signal, the converting means are further adapted to convert the block of said two subsequent 2-bit source words into a corresponding block of two subsequent 3-bit channel words, such that the conversion for said block of two subsequent 2-bit source words is parity preserving.

The measure to convert one (say: the second one) of two subsequent source words into a 3-bit word not identical to the four channel words CW_1 to CW_4 , offers the possibility to detect at the receiver side that a situation existed that encoding of single source words into corresponding single channel words would have led to a violation of the $d=1$ constraint. The encoder now encodes a block of two 2-bit source words into a block of 2 3-bit channel words, such that the encoding of the block is parity preserving, whilst the $d=1$ constraint is satisfied as well.

To embody the encoding of blocks of two 2-bit source words, the device in accordance with the invention may be characterized in that the converting means are adapted to convert the blocks of two consecutive 2-bit source words into the blocks of two consecutive 3-bit channel words in accordance with the coding given in the following table:

block of 2 source words	block of 2 channel words
01 01	100 010
01 00	101 010
11 01	000 010
11 00	001 010

The device in accordance with the invention, for generating a (d,k) sequence, wherein k has a value larger than 5, the device being further provided with means for detecting the position in the bitstream of the source signal where encoding of single 2-bit source words into single 3-bit channel words would lead to a violation of the k -constraint and for supplying a second control signal in response to said detection, may be further characterized in that, in the presence of the second control signal, occurring during the conversion of three consecutive 2-bit source words, the converting means are adapted to convert a block of said three consecutive 2-bit source words into a block of corresponding three consecutive 3-bit channel words,

such that the conversion for said block of three 2-bit source words is parity inverting, the converting means are further adapted to convert two of the three source words in the block into corresponding 3-bit channel words not identical to the four channel words CW_1 to CW_4 , in order to preserve the k constraint.

- 5 This measure enables an encoding of a block of three 2-bit source words into a block of three 3-bit channel words so as to satisfy the k-constraint, and such that the encoding is still parity inverting.

The measure to convert two (say: the second and the third one) of three subsequent source words into a 3-bit word not identical to the four channel words CW_1 to CW_4 , offers the
 10 possibility to detect at the receiver side that a situation existed that the encoding of single 2-bit source words into corresponding single 3-bit channel words would have led to a violation of the k constraint. Upon detection, the decoder is capable of decoding the block of three 3-bit channel words into the corresponding block of three 2-bit source words in the inverse way, as upon encoding.

- 15 To embody the encoding of blocks of three 2-bit source words, the device in accordance with the invention may be characterized in that the converting means are adapted to convert blocks of three consecutive 2-bit source words into blocks of three consecutive 3-bit channel words in accordance with the coding given in the following table:

block of 3 source words	block of 3 channel words
10 10 10	000 010 010
10 10 11	001 010 010
00 10 11	101 010 010
00 10 10	100 010 010

- 20 A device for decoding a stream of data bits of a binary channel signal into a stream of databits of a binary source signal, wherein the bitstream of the channel signal is divided into m-bit channel words, which device comprises deconverting means adapted to deconvert m-bit channel words into corresponding n-bit source words; is characterized in that,
 25 the deconverting means are adapted to deconvert a block of p consecutive m-bit channel words into a corresponding block of p consecutive n-bit source words, such that the conversion for each block of p consecutive m-bit channel words is parity inverting, where n, m and p are integers, $m > n$, $p \geq 1$, and where p is an odd integer and can vary.

It should be noted that published European patent application 199.088A2 discloses a converter for converting n-bit source words into a channel signal in the form of a sequence of m-bit channel words, which channel signal is DC free. The conversion is however not parity inverting.

The invention will be further described in the following figure description, in which

figure 1 shows a first,
figure 2a shows a second,
figure 2b a third, and
figure 3 shows a fourth embodiment of the device,
figure 4 the application of the device in an arrangement for inserting one bit on equidistant positions in the serial source signal, and
figure 5 an embodiment of the decoding device.

Figure 1 shows a device having an input terminal 1, for receiving a stream of databits of a binary source signal S. The terminal 1 is coupled to an input of a shift register 2 having two cells X_1 and X_2 so as to receive two consecutive source bits of the source signal S. The shift register 2 functions as a serial-parallel converter, so as to obtain consecutive 2-bit source words SW. The outputs of the two cells are coupled to two inputs i_1, i_2 of a logic circuit LC, for supplying the logic values (x_1, x_2) of the source bits present in the cells to the logic circuit LC.

The device further includes a second shift register 4 having three cells Y_1, Y_2 and Y_3 . Outputs o_1, o_2 and o_3 of the logic circuit LC are coupled to inputs of the three cells Y_1, Y_2 and Y_3 respectively of the shift register 4, for supplying the logic values (y_1, y_2, y_3) of the channel words. An output 6 of the shift register 4 is coupled to an output terminal 8. The shift register 4 functions as a parallel-serial converter, so as to convert the 3-bit channel words CW supplied by the logic circuit LC into a serial stream of databits of a binary channel signal C.

The logic circuit LC is adapted to convert consecutive 2-bit source words SW into 3-bit channel words, such that the conversion for each 2-bit source word is parity inverting.

That means that the number of 'ones' in the source word to be converted is the 'inverse' of the number of 'ones' in the corresponding channel word, if necessary, after having carried out a

modulo-2 addition on the 'ones' in the channel word. Or, otherwise said: if the number of 'ones' in the source word is even, the number of 'ones' in the channel word will be odd. And: if the number of 'ones' in the source word is odd, the number of 'ones' in the channel word will be even.

- 5 As an example, the converting means LC is adapted to convert the 2-bit source words SW into 3-bit channel words CW in accordance with the following table:

TABLE I

source word (x_1, x_2)		channel word (y_1, y_2, y_3)	
SW ₁	01	CW ₁	101
SW ₂	00	CW ₂	100
SW ₃	11	CW ₃	001
SW ₄	10	CW ₄	000

- 10 It should be noted here, that the first bit in the source word is applied first to the shift register 2 and that the first bit in the channel word is supplied first from the output 6 of the shift register 4.

The bitstream of the channel words is in NRZI (non-return to zero-inverse) notation, which means that a 'one' results in a transition in the write current for recording the channel signal on a magnetic record carrier.

- 15 The device of figure 1 can be used to generate a channel signal C in the form of a (d,k) sequence satisfying the $d=1$ constraint. That means that at least one 'zero' is present between two subsequent 'ones' in the serial datastream of the channel signal. That is, that a concatenation of two or more 'ones' in the channel signal is prohibited.

- 20 It might occur that the unmodified conversion, such as by means of the device of figure 1, of combinations of two subsequent 2-bit source words might violate the $d=1$ constraint. Those combinations are the combinations; '01 01', which by unmodified conversion would lead to the two 3-bit channel words '101 101'; '01 00', which by unmodified conversion would lead to the two 3-bit channel words '101 100'; '11 01', which by unmodified conversion would lead to the two 3-bit channel words '001 101' and '11 00', which by unmodified
25 conversion would lead to the two 3-bit channel words '001 100'.

The occurrence of such combinations should be detected so that a modified encoding of blocks of two 2-bit source words into blocks of two 3-bit channel words can take place. A modified embodiment of a device of figure 1 which is, in addition to the 'normal' encoding of 2-bit source words into 3-bit channel words, capable of detecting the above

identified combinations, and is capable of realizing a modified encoding, such that the $d=1$ constraint in the channel signal is still satisfied, is shown in figure 2a.

The device of figure 2a includes a shift register having four cells X_1 to X_4 so as to receive four consecutive bits (x_1, x_2, x_3, x_4) of the serial bitstream of the source signal S.

5 Outputs of the four cells are coupled to corresponding inputs i_1 to i_4 respectively of the logic circuit LC'. The device further comprises detector unit D1. The detector unit D1 is adapted to detect the position in the serial bitstream of the source signal where unmodified encoding of single source words in the bitstream into corresponding single channel words would lead to a violation of the $d=1$ constraint in the channel signal C, and are adapted to supply a control

10 signal at its output 10 in response to such detection.

The output 10 of the detector unit D1 is coupled to a control signal input 12 of the logic circuit LC'. The logic circuit LC' has six outputs o_1 to o_6 , which are coupled to inputs of cells Y_1 to Y_6 respectively of second shift register 4'.

In the absence of a control signal at the control signal input 12, the logic circuit

15 LC' converts the first 2-bit source word ' $x_1 \cdot x_2$ ' into the three bit channel word ' $y_1 y_2 y_3$ ' in conformity with table I given above. As soon as the detector circuit D1 detects a combination of two 2-bit source words (x_1, x_2, x_3, x_4) which equals one of the combinations given above, the logic circuit LC' converts the combination in accordance with the modified coding as given in the following table:

20

TABLE II

source words	unmodified coding	modified coding
01 01	101 101	100 010
01 00	101 100	101 010
11 01	001 101	000 010
11 00	001 100	001 010

As can be seen from the table, unmodified conversion of the single two 2-bit source words leads to a violation of the $d=1$ constraint, as two 'ones' occur at the boundary

25 between the two channel words obtained. The logic circuit LC' is therefore adapted to convert in a modified coding mode, the blocks of two 2-bit source words given in the left column of the above table into the blocks of two 3-bit channel words as given in the right column in the above table II. As can be seen, no violation of the $d=1$ constraint occurs anymore. Moreover, the modified encoding in the same way is parity preserving. This is correct, for the reason that

twice a parity inverting conversion of a 2-bit source word into a 3-bit channel word, leads to a parity preservation for the combined conversion.

This means in the present situation that, if the number of 'ones' in the blocks of two 2-bit source words is odd (even), the number of 'ones' in the block of two 3-bit channel words obtained is odd (even). Further, one of the two 2-bit source words, which is in the above table the second one, is encoded into a 3-bit channel word which is unequal to one of the four channel words of table I. The reason for this is that on the receiver side, a detection of this 3-bit channel word not belonging to the set of four 3-bit channel words of the table I is possible, so that a corresponding decoding, which is the inverse of the encoding as defined with reference to table II, can be realized.

The block of two 3-bit channel words obtained by means of the encoding in conformity with table II, is supplied by the logic circuit LC' to its outputs o_1 to o_6 , which channel words are supplied to the six cells Y_1 to Y_6 of the shift register 4'. It is clear from the embodiment described that the situations where a modified encoding is needed is detected by means of the detector D1 using the source words.

A different construction of the device for carrying out the modified encoding described with reference to the table II is shown in figure 2b. In this case, detection of the situations where a modified coding should be carried out is decided using the converted channel words. The device of figure 2b includes a detector D1' having 6 inputs for receiving two subsequent 3-bit channel words obtained by means of the unmodified encoding. The detector D1' detects whether the two subsequent 3-bit channel words obtained using the unmodified coding equal one of the four 6-bit sequences given in the middle column under 'unmodified coding' of table II. If so, the detector D1' issues a switching signal at its output 10 and an address signal AD at its output 10'. The switching signal is applied to a switching signal input 45 of the shift register 4". The address signal AD is applied to an address signal input 46 of a ROM 47. The detector D1' generates one of four possible address signals AD1 to AD4, in response to the detection of a corresponding one of the four 6-bit sequences in the middle column of table II. As an example, the address signal AD1 is generated when the detector D1' detects the sequence '101101' and generates the address signal AD4 upon detection of the 6-bit sequence '001100'. The ROM 47 has the 6-bit sequences shown in the right column of table II stored. Upon the receipt of the address signal AD1, the ROM supplies the 6-bit sequence '100 010' at its outputs o_1 to o_6 , and upon the receipt of the address signal AD2, the ROM supplies the 6-bit sequence '101 010' at its outputs. Upon the receipt of the address signal AD3, the ROM supplies the 6-bit sequence '000 010' at its outputs, and upon the

receipt of the address signal AD4, the ROM supplies the 6-bit sequence '001 010' at its outputs. Each memory location of the shift register 4" has now two inputs, one of them being coupled with a corresponding output of the logic circuit LC', the other being coupled to a corresponding output of the ROM 47.

5 In the normal situation, when the $d=1$ constraint is not violated, unmodified conversion is carried out, and the switching signal is absent so that the shift register accepts the bits supplied by the logic circuit LC' via the upper inputs of the shift register 4". If the $d=1$ constraint is violated, the switching signal applied to the switching signal input 45 results in -- the shift register to accept the 6-bit sequence, which is the modified sequence, applied by the
10 ROM to the lower inputs of the shift register 4".

The k -constraint in a (d,k) sequence means that a concatenation of at most k 'zeroes' between two subsequent 'ones' in the channel signal are allowed.

It might occur that the unmodified conversion of three subsequent 2-bit source words might violate the k -constraint:

15 As an example: the sequence of source words '10 10 10' would by unmodified conversion lead to the three 3-bit channel words '000 000 000'. If a (d,k) sequence should be obtained where k equals 6, 7 or 8, such combination of three 3-bit channel words should not occur.

20 Another example is the sequence of source words '10 10 11' which by unmodified conversion would lead to the three 3-bit channel words '000 000 001'. This combination of three 3-bit channel words does not satisfy a $k=6$ or $k=7$ constraint. Moreover, this combination of three 3-bit channel words can follow a previous channel word that ends with a '0', so that it might lead to a violation of a $k=8$ constraint. Further, the combination ends with a '1', so that it might lead to a violation of the $d=1$ constraint, if the combination is
25 followed by a 3-bit channel word that starts with a '1'. An equivalent reasoning is valid for the sequence of source words '00 10 10'.

A further example is the sequence of source words '00 10 11' which by unmodified conversion would lead to the three 3-bit channel words '100 000 001'. This combination can, in the same way as given above, lead to a violation of the $d=1$ constraint.

30 The occurrence of such combinations should be detected so that a modified encoding can take place. An embodiment of a device which is, in addition to the 'normal' encoding of 2-bit source words into 3-bit channel words, capable of detecting the above identified combinations, and is capable of realizing a modified encoding, is shown in figure 3.

The device of figure 3 includes a shift register 2" having six cells X_1 to X_6 so as to receive six consecutive bits of the serial bitstream of the source signal S. Outputs of the six cells are coupled to corresponding inputs i_1 to i_6 respectively of the logic circuit LC". The device further comprises detector means D2. The detector means D2 are adapted to detect the position in the serial bitstream of the source signal where unmodified encoding of the bitstream would lead to a violation of the k-constraint in the channel signal C, and are adapted to supply a control signal at its output 15 in response to such detection.

The outputs of the six cells are also coupled to four inputs i_1 to i_6 respectively of logic circuit LC". The output 15 of the detector means D2 is coupled to a control signal input 16 of the logic circuit LC". The logic circuit LC" has nine outputs o_1 to o_9 , which are coupled to inputs of cells Y_1 to Y_9 respectively of second shift register 4".

In the absence of a control signals at the control signal inputs 12 and 16, the logic circuit LC' converts a single 2-bit source word ' $x_1 x_2$ ' into a single 3-bit channel word ' $y_1 y_2 y_3$ ' in conformity with table I given above. As soon as the detector circuit D1 detects a block of two 2-bit source words ' $x_1 x_2, x_3 x_4$ ' which equals one of the combinations given in table II above, the logic circuit LC" converts the combination in accordance with the conversion rule as given in table II, so as to obtain a block of two 3-bit channel words ' $y_1 y_2 y_3 y_4 y_5 y_6$ '.

As soon as the detector D2 detects a block of three 2-bit source words ' $x_1 x_2 x_3 x_4, x_5 x_6$ ' which equals one of the combinations given above, the logic circuit LC" converts the block in accordance with the modified coding as given in the following table, so as to obtain a block of three 3-bit channel words:

TABLE III

source words	unmodified coding	modified coding
10 10 10	000 000 000	000 010 010
10 10 11	000 000 001	001 010 010
00 10 11	100 000 001	101 010 010
00 10 10	100 000 000	100 010 010

The logic circuit LC" is adapted to convert in a modified coding mode, the blocks of three 2-bit source words given in the left column of the above table III into the blocks of three 3-bit channel words as given in the right column in the above table. By realizing the modified encoding as per table III, a channel signal has been obtained which satisfies the $k=8$ constraint. Moreover, the modified encoding in the same way is parity inverting. This means in the

present situation that, if the number of 'ones' in the combination of three 2-bit source words is odd (even), the number of 'ones' in the combination of the three 3-bit channel words obtained is even (odd). Further, two of the three 2-bit source words, which is in the above table the second one and the third one, is encoded into a 3-bit channel word which is unequal to one of the four channel words of table I. The reason for this is that on the receiver side, a detection of these two consecutive 3-bit channel words not belonging to the set of four 3-bit channel words of the table I is possible, so that a corresponding decoding, which is the inverse of the encoding as defined with reference to table III, can be realized.

The combination of three 3-bit channel words obtained by means of the encoding in conformity with table III, is supplied by the logic circuit LC" to its outputs o_1 to o_9 , which channel words are supplied to the nine cells Y_1 to Y_9 of the shift register 4". The serial datastream of the channel signal C is supplied to the output terminal 8.

It will be evident that, in the same way as described with reference to figure 2b, the detection of the violation of the k-constraint can be done on the channel signal level, instead of the source signal level.

It has been said previously that other conversion rules for converting single 2-bit source words into single 3-bit channel words are possible. Those conversion rules are given in the following three tables.

TABLE IV

source word (x_1, x_2)		channel word (y_1, y_2, y_3)	
SW ₁	01	CW ₁	101
SW ₂	00	CW ₂	001
SW ₃	11	CW ₃	100
SW ₄	10	CW ₄	000

TABLE V

source word (x_1, x_2)		channel word (y_1, y_2, y_3)	
SW ₁	01	CW ₁	000
SW ₂	00	CW ₂	100
SW ₃	11	CW ₃	001
SW ₄	10	CW ₄	101

TABLE VI

source word (x_1, x_2)		channel word (y_1, y_2, y_3)	
SW ₁	01	CW ₁	000
SW ₂	00	CW ₂	001
SW ₃	11	CW ₃	100
SW ₄	10	CW ₄	101

It is evident that extensions of those conversion rules for encoding blocks of two or three 2-bit source words into blocks of two or three 3-bit channel words can be obtained using the teachings given above.

A further embodiment of an encoder is explained with reference to the following table VII. This table shows the conversion rule for an encoder capable of encoding 3-bit source words into 4-bit channel words.

TABLE VII

source word (x_1, x_2, x_3)		channel word (y_1, y_2, y_3, y_4)	
SW ₁	001	CW ₁	0000
SW ₂	000	CW ₂	0001
SW ₃	011	CW ₃	0100
SW ₄	010	CW ₄	0101
SW ₅	101	CW ₅	1000
SW ₆	100	CW ₆	1001
SW ₇	111	CW ₇	1010
SW ₈	110	CW ₈	0010

As has been said previously, the devices described above are very suitable in combination with a converter unit in which one bit is inserted after each q bits in a serial datastream in order to realize a polarity conversion, or not. Figure 4 shows such combination, where the converter unit 40 is followed by the device 7' in accordance with the present invention 41, which device 7' is subsequently followed by a 1T-precoder 42, well known in the art. The output signal of the 1T-precoder 42 is applied to a control signal generator 43, which generates the control signal for the converter unit 40, so as to control whether a '0' or a '1' is inserted in the serial datastream applied to the device 7'. Inserting a '0' or a '1' bit always

leads to either an increase and decrease, respectively, or vice versa, in the running digital sum value at the output of the precoder 42.

By means of the arrangement shown in figure 4 it is possible to embed a tracking tone of a certain frequency in the serial datastream, or keep the DC content of the datastream to zero.

Further, when the device 7' is adapted to generate a (d,k) sequence as explained above, it causes the output signal of the arrangement of figure 4 to be a (d,k) RLL output signal.

Embodiments of the converter 40 are given in Bell System Technical Journal, Vol 53, No. 6, pp. 1103-1106.

Figure 5 shows a decoding device for decoding the serial datastream obtained by the encoding device of figure 3 so as to obtain a binary source signal. The decoding device has an input terminal 50 for receiving the channel signal, which input terminal 50 is coupled to an input 56 of a shift register 51, comprising nine cells Y_1 to Y_9 . The shift register 51 functions as a serial-parallel converter so that blocks of three 3-bit channel words are applied to inputs i_1 to i_9 of a logic circuit 52. The logic circuit 52 comprises the three tables I, II and III. Outputs o_1 to o_6 of the logic circuit 52 are coupled to inputs of cells X_1 to X_6 of a shift register 54, which has an output 57 coupled to an output terminal 55. A detector circuit 53 is present having inputs i_1 to i_6 coupled to outputs of cells Y_4 to Y_9 respectively of the shift register 51, and outputs o_1 and o_2 coupled to control inputs c_1 and c_2 respectively of the logic circuit 52. The detector circuit 53 is capable of detecting a '010' bit pattern in the cells Y_4 , Y_5 and Y_6 of the shift register 51 and is capable of detecting a bit pattern '010010' in the cells Y_4 to Y_9 of the shift register 51.

Upon detection of the '010010' bitpattern, the detector circuit 53 generates a control signal on its output o_2 , and upon detection of a '010' bit pattern in the cells Y_4 , Y_5 and Y_6 , whilst there is no '010' bit pattern in the cells Y_7 , Y_8 and Y_9 , it generates a control signal on its output o_1 .

In the absence of the control signals, the logic circuit 52 converts the 3-bit channel word stored in the cells Y_1 , Y_2 and Y_3 into its corresponding 2-bit source word, as per the conversion table I, and supplies the 2-bit source word to the cells X_1 and X_2 . In the presence of the control signal at the input c_1 , the logic circuit 52 converts the block of two 3-bit channel words stored in the cells Y_1 to Y_6 into a block of two 2-bit source words, as per the conversion table II, and supplies the two 2-bit source words to the cells X_1 to X_4 . In the presence of the control signal at the input c_2 , the logic circuit 52 converts the block of three 3-bit channel words stored in the cells Y_1 to Y_9 into a block of three 2-bit source words, as per the conversion table III, and supplies the three 2-bit source words to the cells X_1 to X_6 . In this

way, the serial datastream of the channel signal is converted into the serial datastream of the source signal.

The encoded information supplied to the input 50 could have been obtained from reproducing the information from a record carrier, such as a magnetic record carrier 23 or an optical record carrier 23'. The device of figure 5 thereto comprises a read unit 62 for reading the information from a track on the record carrier, where the unit 62 comprises a read head 64 for reading the information from said track.

Next, another parity inverting 2-to-3 bit conversion code will be described, resulting in a (1,7) sequence. The main conversion table is as follows:

source word		Channel word	
SW ₁	01	CW ₁	x0x
SW ₂	00	CW ₂	001
SW ₃	11	CW ₃	010

In this table, conversion of the source word '01' is dependent of the last bit of the channel word obtained from converting the directly preceding 2-bit source word. When, this last bit is a '0' bit, the conversion results into the 3-bit word '101' and when this last bit is a '1' bit, the conversion results into the 3-bit word '000'.

A first substitution table is present for converting specific blocks of two 2-bit source words. This first substitution table is as follows:

block of 2 source words	block of 2 channel words
010 01	010 100
010 00	010 000
010 11	000 100

A second substitution table is present for converting specific blocks of three 2-bit source words. This second substitution table is as follows:

block of 3 source words	block of 3 channel words
10 10 01	000 100 100
10 10 00	001 100 000
10 10 11	010 100 100
10 10 10	010 100 000

A third substitution table is present for converting specific blocks of four 2-bit source words.

This third substitution table is as follows:

5

block of 4 source words	block of 4 channel words
10 10 00 10	000 100 100 100
10 10 10 10	010 100 100 100

Further, unmodified conversion of the following sequence 01 11 01 xy could lead to the following channel sequence: 101 010 101 010, where the conversion of the first 2-bit source word apparently led to the 3-bit channel word '101' and xy is a 2-bit source word leading to the 3-bit channel word '010'. Such sequence is unwanted, as it violates a requirement for the length of a repeated minimum transition runlength (RMTR). Therefore, upon occurrence of such sequence, this sequence is converted into the sequence 001 000 000 010.

Whilst the invention has been described with reference to preferred embodiments thereof, it is to be understood that these are not limitative examples. Thus, various modifications may become apparent to those skilled in the art, without departing from the scope of the invention, as defined by the claims. As an example, the decoding device of figure 5 could be modified into a device in which the detector 53 detects the various modified decoding situations from the decoded information, instead from the encoded information, as disclosed in figure 5. It should further be noted that, as an example, the converter unit 7' and the precoder 42 could have been into one combined unit, where, dependent of incoming n-bit source words, via a conversion table, those n-bit source words are directly converted into 3-bit output words of the combined unit. Further, it should be noted that the parity inverting conversion as claimed could have been obtained by applying a parity preserving coder, such as described in USP 5,477,222, and EXORing the 2-bit source words with either '10' or '01'

(EXORing, in the sense of: msb of the 2-bit source word with the msb of '10' or '01', and lsb of the 2-bit source word with the lsb of '10' or '01'), prior to applying the 2-bit source words to the parity preserving coder.

Further, any reference signs do not limit the scope of the claims. The invention can be

- 5 implemented by means of both hardware and software, and several "means" may be represented by the same item of hardware. The word 'comprising' does not exclude the presence of other elements or steps than those listed in a claim. Also, the word "a" or "an" preceding an element does not exclude the presence of a plurality of such elements. In addition, the invention lies in each and every novel feature or combination of features.

CLAIMS:

1. A device for encoding a stream of databits of a binary source signal into a stream of databits of a binary channel signal, wherein the bitstream of the source signal is divided into n-bit source words, which device comprises converting means adapted to convert said source words into corresponding m-bit channel words, characterized in that, the converting means are adapted to convert a block of p consecutive n-bit source words into a corresponding block of p consecutive m-bit channel words, such that the conversion for each block of p consecutive n-bit source words is parity inverting, where n, m and p are integers, $m > n \geq 1$, $p \geq 1$, and where p is an odd integer and can vary.

2. A device as claimed in claim 1, characterized in that, $m=n+1$.

3. A device as claimed in claim 2, characterized in that, $n=2$.

4. Device as claimed in claim 3, characterized in that, the device is adapted to convert single source words into corresponding single channel words in accordance with the following table:

source word		Channel word	
SW ₁	00	CW ₁	100
SW ₂	01	CW ₂	101
SW ₃	10	CW ₃	000
SW ₄	11	CW ₄	001

5. A device as claimed in claim 3 or 4, wherein the converting means are adapted to convert 2-bit source words into corresponding 3-bit channel words, so as to obtain a channel signal in the form of a (d,k) sequence, where $d=1$, the device further comprising means for detecting the position in the bitstream of the source signal where encoding of single 2-bit source words into corresponding single channel words would lead to a violation of the d-constraint at the channel word boundaries and for supplying a control signal in response to

said detection, characterized in that, in the absence of the control signal, the converting means are adapted to convert single 2-bit source words into corresponding single 3-bit channel words, such that the conversion for each 2-bit source word is parity inverting.

- 5 6. Device as claimed in claim 5, wherein, in the presence of the control signal, occurring during the conversion of two consecutive source words, the converting means are adapted to convert a block of said two consecutive 2-bit source words into a block of two corresponding 3-bit channel words, such that one of the two source words in the block of source words is converted into a 3-bit channel word which is not identical to one of the four
10 channel words CW_1 to CW_4 , in order to preserve the $d=1$ constraint, characterized in that, in the presence of said control signal, the converting means are further adapted to convert the block of said two subsequent 2-bit source words into a corresponding block of two subsequent 3-bit channel words, such that the conversion for said block of two subsequent 2-bit source words is parity preserving.

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7. Device as claimed in claim 1 or 6, characterized in that, the converting means are adapted to convert blocks of two consecutive 2-bit source words into blocks of two consecutive 3-bit channel words in accordance with the coding given in the following table:

block of 2 source words	block of 2 channel words
01 01	100 010
01 00	101 010
11 01	000 010
11 00	001 010

20

8. Device as claimed in claim 1, 6 or 7, where k has a value larger than 5, the device being further provided with means for detecting the position in the bitstream of the source signal where encoding of single 2-bit source words into single 3-bit channel words would lead to a violation of the k -constraint and for supplying a second control signal in
25 response to said detection, characterized in that, in the presence of the second control signal, occurring during the conversion of three consecutive 2-bit source words, the converting means are adapted to convert a block of said three consecutive 2-bit source words into a block of corresponding three consecutive 3-bit channel words, such that the conversion for said block of three 2-bit source words is parity inverting, the converting means are further adapted to

convert two of the three source words in the block into corresponding 3-bit channel words not identical to the four channel words CW_1 to CW_4 , in order to preserve the k constraint.

9. Device as claimed in claim 1 or 8, characterized in that, the converting means are adapted to convert blocks of three consecutive 2-bit source words into blocks of three consecutive 3-bit channel words in accordance with the coding given in the following table:

block of 3 source words	block of 3 channel words
10 10 10	000 010 010
10 10 11	001 010 010
00 10 11	101 010 010
00 10 10	100 010 010

10. Device as claimed in anyone of the preceding claims, characterized in that the conversion means are adapted to carry out a signal processing on the binary source signal equivalent to the conversion of the blocks of p consecutive source words into the blocks of p consecutive channel words, followed by an aT precoding of said channel words.

11. Device as claimed in claim 1 or 10, characterized in that it further comprises bitadding means for adding one bit to subsequent blocks of q bits of the source signal.

12. Device as claimed in anyone of the preceding claims, characterized in that it further comprises means for recording the stream of databits of the binary channel signal in a track on the record carrier.

13. Method of encoding a stream of databits of a binary source signal into a stream of databits of a binary channel signal, wherein the bitstream of the source signal is divided into n -bit source words, the method comprising the step of converting said source words into corresponding m -bit channel words, characterized in that, the converting step comprises the conversion of a block of p consecutive n -bit source words into a corresponding block of p consecutive m -bit channel words, such that the conversion for each block of p consecutive n -

bit source words is parity inverting, where n , m and p are integers, $m > n \geq 1$, $p \geq 1$, and where p can vary.

14. Device for decoding a stream of data bits of a binary channel signal into a
5 stream of databits of a binary source signal, wherein the bitstream of the channel signal is
divided into m -bit channel words, which device comprises deconverting means adapted to
deconvert m -bit channel words into corresponding n -bit source words, characterized in that,
the deconverting means are adapted to deconvert a block of p consecutive m -bit channel words
into a corresponding block of p consecutive n -bit source words, such that the conversion for
10 each block of p consecutive m -bit channel words is parity inverting, where n , m and p are
integers, $m > n$, $p \geq 1$, and where p is an odd integer and can vary.

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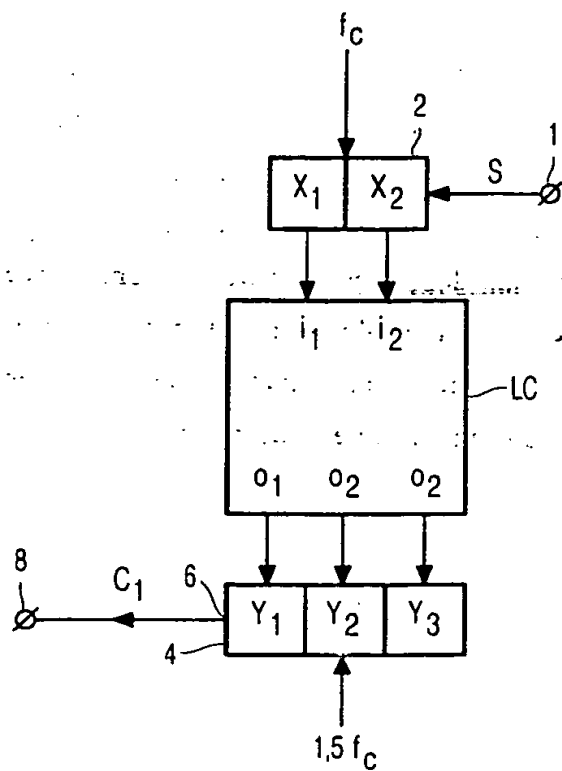


FIG. 1

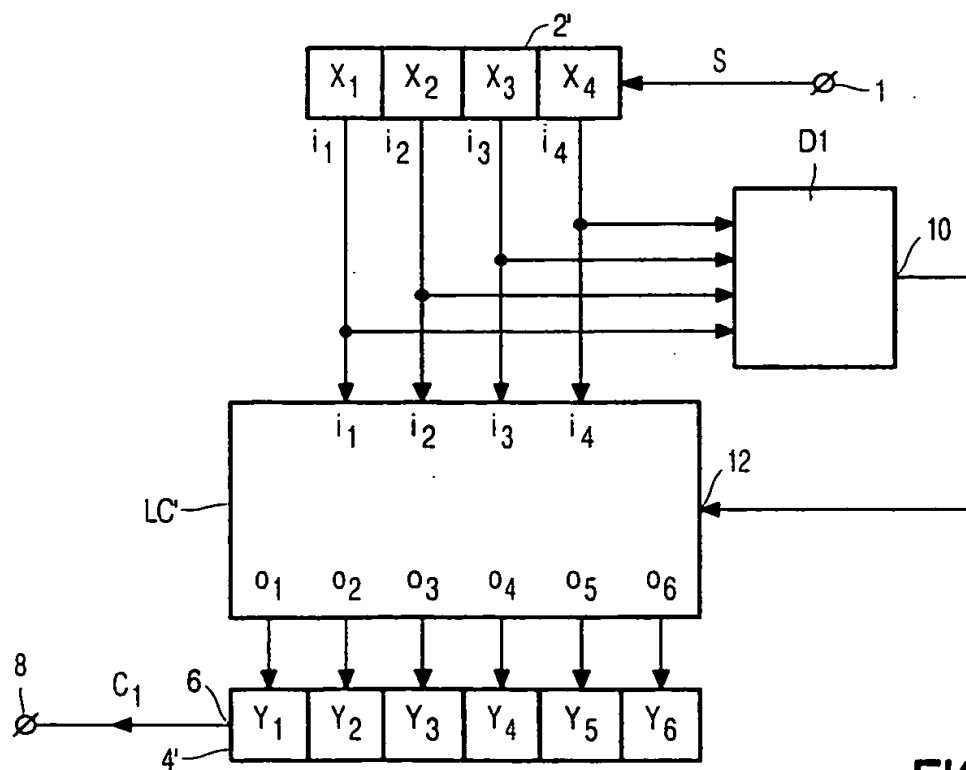
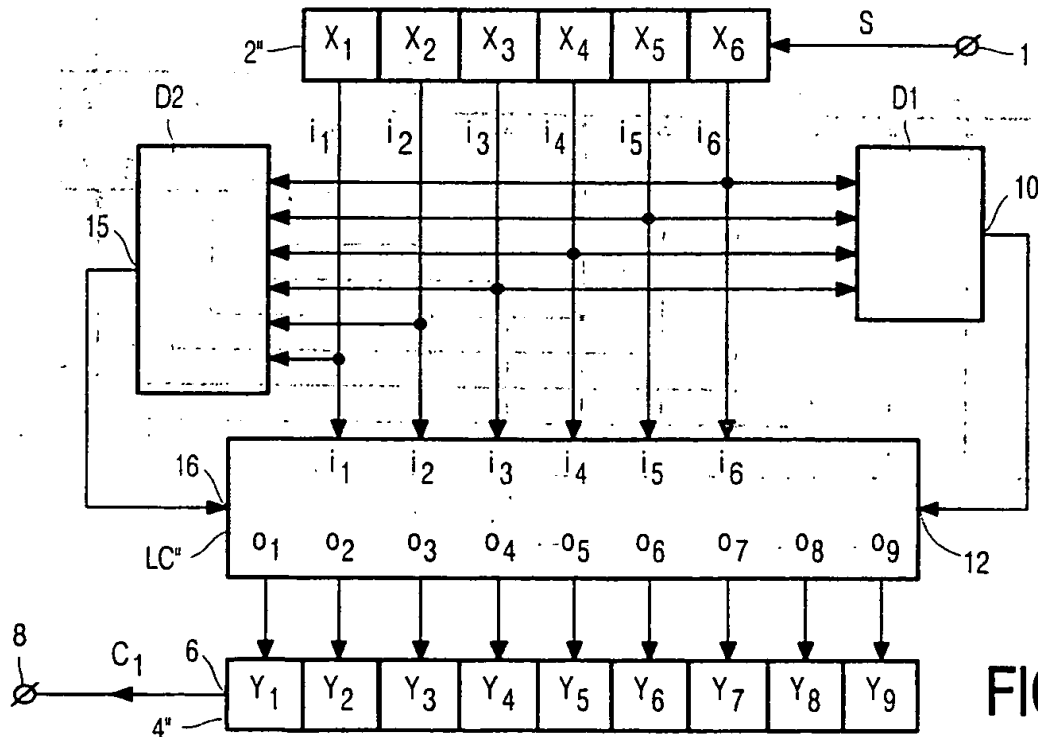
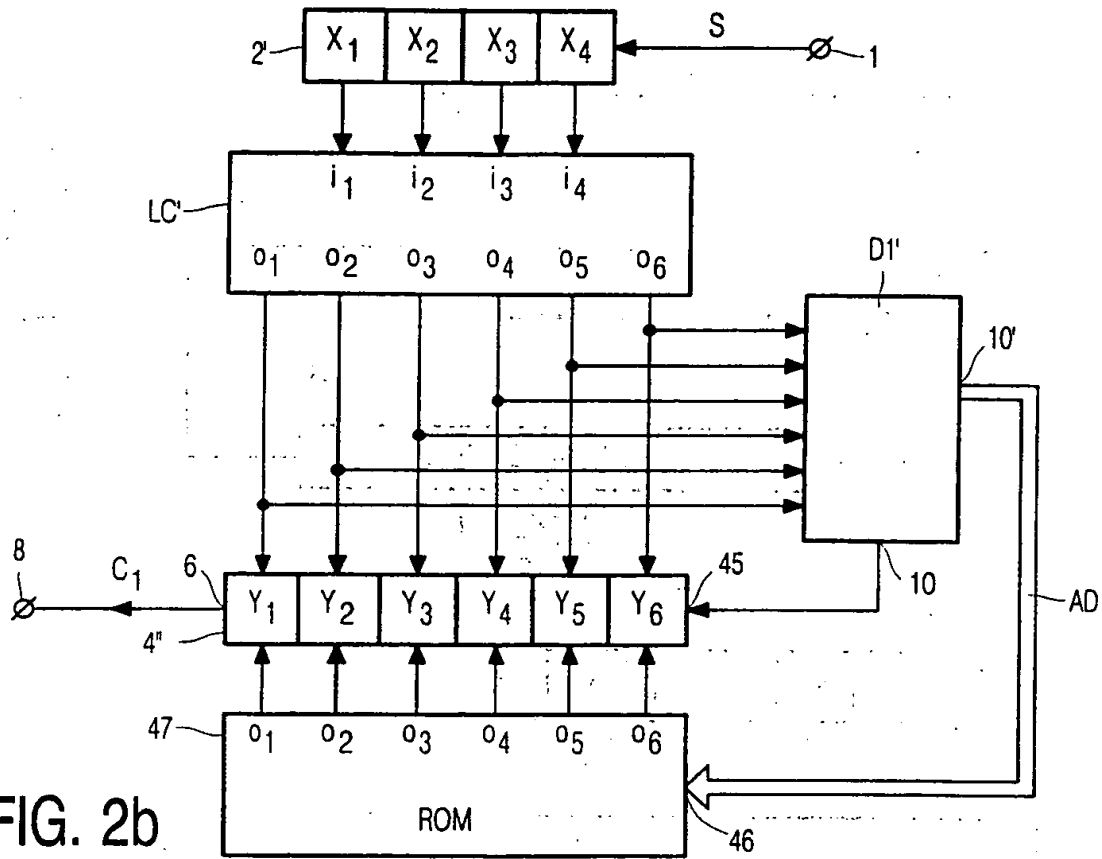


FIG. 2a

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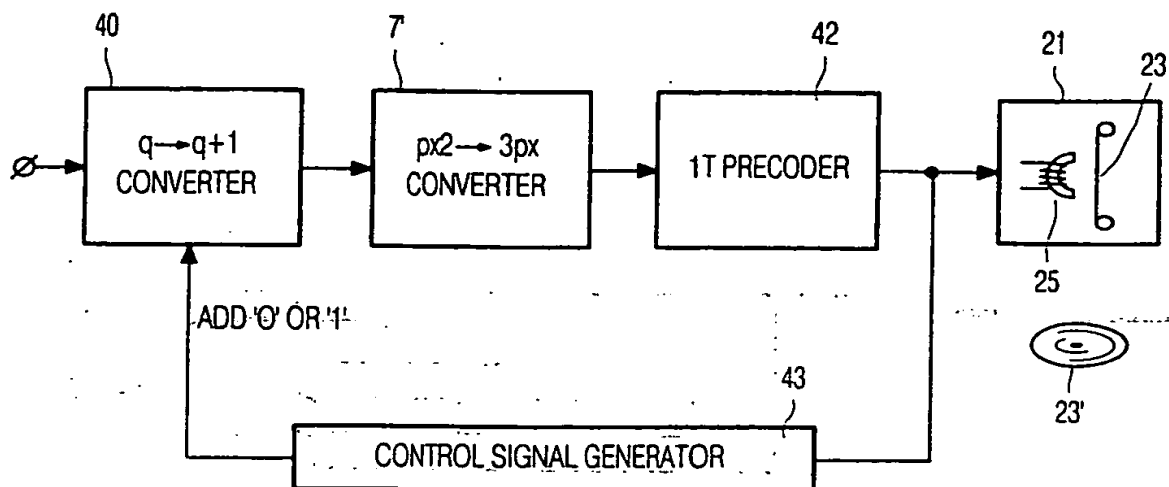


FIG. 4

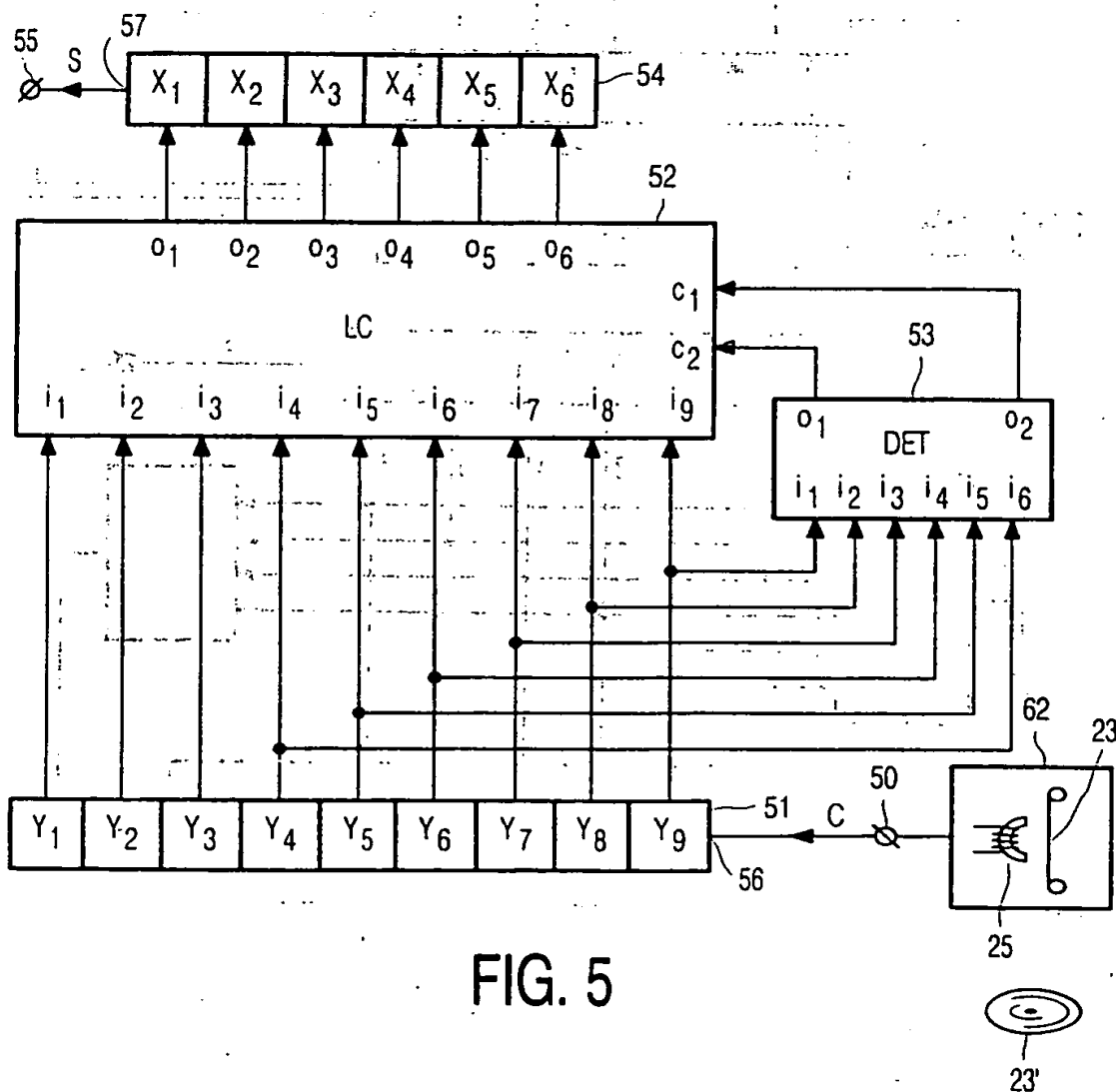


FIG. 5

INTERNATIONAL SEARCH REPORT

International Application No

PCT/EP 00/01340

A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 H03M5/14 G11B20/14

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 7 H03M G11B

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched.

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	US 5 477 222 A (KAHLMAN JOSEPHUS A H M ET AL) 19 December 1995 (1995-12-19) cited in the application the whole document	1-3,5, 10-14
A	---	4,6-9
Y	US 5 278 902 A (NUGENT STEVEN F) 11 January 1994 (1994-01-11) figures 5C,7A column 6, line 15 - line 25 column 7, line 1 - line 3 column 7, line 18 - line 22 ---	1-3,5, 10-14
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30 May 2000

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INTERNATIONAL SEARCH REPORT

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C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	<p>SIEGEL P H ET AL: "MODULATION AND CODING FOR INFORMATION STORAGE"</p> <p>IEEE COMMUNICATIONS MAGAZINE, US, IEEE SERVICE CENTER, PISCATAWAY, N.J., vol. 29, no. 12, 1 December 1991 (1991-12-01), pages 68-86, XP000287983</p> <p>ISSN: 0163-6804</p> <p>page 74 -page 75</p> <p>page 83 -page 84</p> <p>-----</p>	1-14

INTERNATIONAL SEARCH REPORT

information on patent family members

International Application No

PCT/EP 00/01340

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
US 5477222 A	19-12-1995	EP 0624000 A JP 7050588 A	09-11-1994 21-02-1995
US 5278902 A	11-01-1994	NONE	

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